

END SEMESTER EXAMINATION, JUNE 2019
MCA THIRD SEMESTER

COURSE CODE: MCA206

COURSE TITLE: THEORY OF AUTOMATA

[Time allotted: Three hours]**[Max. Marks: 100]****Note:** Attempt all Sections & Questions.

Section (A)

Q. 1. Attempt all questions.**(2 x 10 = 20)**

- a. Design a DFA that will accept a unary string divisible by 5.
- b. Write short note on regular grammar.
- c. Write short note on Moore & Mealy Machine.
- d. A Language for which no DFA exist is a _____ (Choose any one).
 - i. Regular Language
 - ii. Non- Regular Language
 - iii. May be Regular
 - iv. Cannot be said
- e. A push down automata can be represented as $PDA = \epsilon\text{-NFA} + [\text{stack}]$. (True or false).
- f. Design a DFA that will accept a string of even number of 0's.
- g. The format: $A \rightarrow aB$ refers to which of the following? (Choose any one).
 - i. Chomsky Normal Form
 - ii. Greibach Normal Form
 - iii. Backus Naur Form
 - iv. None of the mentioned
- h. Which of the following are always unambiguous? (Choose any one)
 - i. Deterministic Context free grammars
 - ii. Non-Deterministic Regular grammars
 - iii. Context sensitive grammar
 - iv. None of the mentioned
- i. State and Proof Arden's Theorem.
- j. Discuss about Chomsky Classification of Grammars in short.

Section (B)

Q. 2. Attempt any 5 questions.**(7 x 5 = 35)**

- a. Construct a Turing Machine to accept the set L of all String over $\{0,1\}$ ending with 111.
- b. Consider the following production:
 $S \rightarrow aB/bA, \quad A \rightarrow aS/bAA/a, \quad B \rightarrow bS/aBB/b$
For the string aaabbabbba, find
 - a. Leftmost derivation
 - b. Rightmost derivation
 - c. Parse tree
- c. Design a DFA which accepts the language
 $L = \{w \in (0,1)^* \mid \text{second symbol of } w \text{ is '1' and fourth symbol of } w \text{ is '1'}\}$.
- d. Design a DFA for the language
 $L = \{w \in (a,b)^* \mid n_a(w) \bmod 3 > 0\}$
- e. Find a regular expression corresponding to each of the following subsets of $\{a,b\}^*$.
 - a. The set of all strings containing exactly 2b's.

- b. The set of all strings containing at least 2b's.
 c. The set of all strings containing a most 2b's.
 d. The set of all strings containing the substring bb.
- f. Construct the transition system for given regular expression: $R = 00 + 1 + 010^*1$
- g. Turing Machine M described by the transition table given in table. Describe the processing of (a) 011 (b) 0011 (c) 001 which of the above strings are accepted by M?

Transition Table					
Present State	Tape Symbol				
	0	1	x	Y	b
$\rightarrow q_1$	xRq ₂				bRq ₅
q ₂	0Rq ₂	yLq ₃		yRq ₂	
q ₃	0Lq ₄		xRq ₅	yLq ₃	
q ₄	0L q ₄		xRq ₁		
q ₅				yRq ₅	bRq ₆
q₆					

Section (C)

Attempt **any 3** questions.

(15 x 3 = 45)

Q. 3.

- a. Define CNF & GNF with rule. Convert given grammar in GNF:
 $E \rightarrow E + T / T$
 $T \rightarrow T * F / F$
 $F \rightarrow (E) / a$

- b. Convert Mealy Machine into Moore Machine:

PRESENT STATE	NEXT STATE			
	Input a		Input b	
	State	Output	State	Output
$\rightarrow q_0$	q ₀	0	q ₁	1
q ₁	q ₁	1	q ₂	0
q ₂	q ₂	0	q ₀	1

Q. 4.

- a. Differentiate between Turing machine & Linear bounded automata. Design a Turing machine to recognize the language $\{a^n b^n c^n / n \geq 1\}$. Check the string "aabbcc" for acceptance.

- b. Consider the following grammar : $S \rightarrow A/bb$ $A \rightarrow B/b$ $B \rightarrow S/a$
identified unit production and remove them.

Q. 5.

- a. Differentiate between DPDA & NPDA. Design a PDA accepting the set of all strings over $\{a,b,c\}$ for $\{a^n b^{n+m} c^m \mid n \geq 1, m \geq 1\}$.
- b. What is an ambiguous grammar? Check that grammar $S \rightarrow aAS/a$, $A \rightarrow SbA/SS/ba$ is an ambiguous grammar or unambiguous grammar.

Q. 6.

- a. Design a PDA for the language $\{ww^R \mid w \in (a,b)^*\}$.
- b. Construct a reduced grammar equivalent to the grammar G:
 $S \rightarrow aAa$, $A \rightarrow Sb/bCC/DaA$, $C \rightarrow abb/DD$, $D \rightarrow aDA$, $E \rightarrow aC$

Q. 7.

- a. Construct a context-free grammar G which accepts a given PDA where PDA moves are:
 $(q_0, b, z_0) = \{(q_0, zz_0)\}$
 $(q_0, \wedge, z_0) = \{(q_0, \wedge)\}$
 $(q_0, b, z) = \{(q_0, zz)\}$
 $(q_0, a, z) = \{(q_1, z)\}$
 $(q_1, b, z) = \{(q_1, \wedge)\}$
 $(q_1, a, z_0) = \{(q_0, z_0)\}$
- b. Convert the following NFA into DFA.

